

# Introduction to Computer Engineering – EECS 203

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## RSE processor

- Already understand building FSMs
- Can use array of latches to store multiple bits: register
- Consider simple processor, called RSE (Rob's simplified example)

## RSE arithmetic instructions

- add  $R_D, R_{S1}, R_{S2}$
- sub  $R_D, R_{S1}, R_{S2}$

Do computation on source registers and put result in destination register

## Branch instructions

- blz  $R_T, R_C$ 
  - Set  $PC$  to  $R_T$  if  $R_C < 0$
- bz  $R_T, R_C$ 
  - Set  $PC$  to  $R_T$  if  $R_C = 0$

## Change in style

- Micro-controller based design
- In this lecture, I want a lot of help and participation
- You now have the fundamental knowledge to design a processor
- Let's build a simple one on paper
- You'll be programming a slightly more complex processor in next week's lab assignment

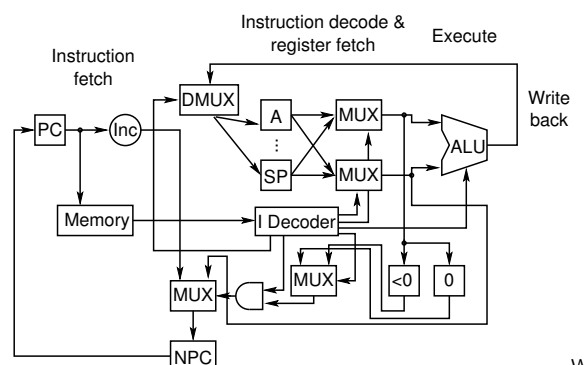
## RSE registers

- All registers are 8-bit
- Four general-purpose registers,  $A, B, C,$  and  $D$ 
  - Used to do computation
- Program counter  $PC$
- Stack pointer  $SP$  (sometimes called  $TOS$  for top of stack), which may also be used as a general-purpose register
- ALU capable of adding (0) and subtracting (1)

## RSE data motion

- ldm  $R_D, [R_S]$ 
  - Load from memory location indicated by the source register into destination register
- stm  $[R_D], R_S$ 
  - Store to memory location indicated by the destination register from source register
- ldi  $R_D, I$ 
  - Load immediate into destination register
- ldpc  $R_S$ 
  - Load from program counter to destination register

## Architecture



What

## Instruction encoding

- How many instructions?
- Worst-case operands?
  - 3 registers (each how many bits?)
  - 1 register and 1 immediate
  - To pack or not to pack?

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## Memory

- Acts like a collection of byte-wide registers
- Address using a decoder
- Can put other devices at some memory locations
  - Memory-mapped input-output
- Can also use special-purpose output instructions or registers
- Let's build some from D flip-flops
- Multiplexing address and data lines?

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## Example high-level code

Sum up the contents of memory locations 2-6

- 1  $A = 0$
- 2 For  $B$  from 2 to 6
- 3  $A = A + [B]$

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## Error conditions

- What happens on overflow or underflow?
- Special register?
- Special value associated with each register?
- Single-instruction compare and branch?
- Advantages and disadvantages of each?

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## Initialization

- Chip has reset line
- Set  $PC$  to byte 2
- Start running...

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## Program counter

Every clock tick the processor

- Fetches an instruction from the memory location pointed to by  $PC$
- Decodes the instruction
- Fetches the operands
- Executes the instruction
- Stores the results
- Increments the program counter
- Can jump to another code location by moving a value into the  $PC$

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## Example low-level code

Sum up the contents of memory locations 2-6

```

2. A = 0                                ldi A, 0 or sub A, A, A
4. B = 2                                ldi B, 2
6. C = [B] (loop start point)          ldm C, [B]
8. A = A + C                             add A, A, C
10. B = B + 1                             ldi C, 1 — add B, B, C
14. C = 6 (loop start)                   ldi C, 6
16. If B ≤ 6 (B < 7) branch to C         ldi D, 7 — sub D, B, D — blz C,
D
(Done)

```

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## Assemble to our encodings

- After assembling, can put program contents into memory, starting at byte 2
- Compiling from higher-level languages also possible

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## Example high-level code

Sum up the contents of memory locations 2-6

- 1  $i = 0$
- 2 For  $j$  from 2 to 6
- 3      $i = i + [j]$

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## Today's topics

- Architecture
- Assembly
- Compilation
- PIC16C74A

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## Computer geek culture references

- Building multicontroller-based devices for the fun of it
- <http://www.bdmicro.com>
- <http://www.commlinx.com.au/microcontroller.htm>
- <http://members.home.nl/bzijlstra/>
- <http://www.robotcafe.com/dir/Companies/Hobby/more3.shtml>
- Etc.

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## Lesson

- With only a few registers and instructions, powerful actions are possible
- Less time and power efficient than special-purpose hardware design
- Instruction processors are flexible
- Allows massive use of a single type of IC
- Assembly is painful
- However, much better than doing hardware design
- Compilation also possible

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## Assigned reading

- M. Morris Mano and Charles R. Kime. *Logic and Computer Design Fundamentals*. Prentice-Hall, NJ, fourth edition, 2008
- Refer to Chapter 7 and 8
- Read Sections 9.1–9.7, 10.1–10.6, 10.8

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