

Homework One

EECS 303: Advanced Digital Logic Design

Assigned 25 September

Due date to be announced in class, definitely not due before 2 October

You may discuss the assignment with your classmates. However, you need to understand and write the solutions independently.

- (10 points)** The Boolean function $f(a_1, a_0, b_1, b_0)$ has two two-bit inputs, for a total of four input bits. $a = (a_1, a_0)$ and $b = (b_1, b_0)$ are two-bit binary numbers. $f(a_1, a_0, b_1, b_0)$ is true if and only if $a = b$.
 - Write a Boolean formula for f .
 - Write the truth table for f .
 - Draw a Karnaugh map for f .
 - a and b are each two-bit numbers. If a and b were each n -bit numbers, what is the minimal number of AND gates that would be required for a SOP implementation of

$$f(a_{n-1}, \dots, a_1, a_0, b_{n-1}, \dots, b_1, b_0)$$

The AND gates may have more than two inputs. Explain your answer.

- (10 points)** Simplify the following Boolean functions using algebraic manipulation, K-Maps, and the Quine-McCluskey algorithm, i.e., solve each problem using each technique. When doing algebraic manipulation, indicate the law you use with each step. For example, if you would like to convert from $a + b\bar{c} + d$ to $a + (\bar{b}\bar{c})d$, you can indicate that you're using the following law: $(a + b) = \bar{a}\bar{b}$. You don't need to state the name of the law as long as you give its canonical form/definition. Treat don't-cares as zeros when doing algebraic minimization.

Do any of these problems have special properties? Which one(s)? What is the property called? What impact does it have on minimization?

(a) $f(a, b, c, d) = \sum(0, 5, 6, 10, 11, 13) + d(4, 8, 14)$

(b) $f(a, b, c) = \sum(0, 2, 3, 4, 5, 7)$

(c) $f(a, b, c, d) = \sum(1, 7, 8, 11) + d(0, 3, 14)$

- (10 points)** Minimize the following function using whatever manual technique you prefer

$$f(a, b, c, d, e) = \sum(0, 1, 2, 5, 6, 8, 10, 11, 15, 17) + d(4, 9, 12, 13, 16, 21, 26, 29, 30, 31)$$

- (5 points)** Minimize the following functions using K-maps. Give your answer in POS form. In this notation, the zeros and don't-cares are specified instead of the ones and don't-cares. You can use either of the techniques presented in class.

(a) $f = \prod(1, 4, 5, 9, 11, 14) + d(0, 2, 7, 8, 13)$

(b) $f = \prod(2, 7, 10, 11, 14) + d(1, 4, 6)$

5. **(5 points)** Define the following terms:
 - (a) Observability Don't-care
 - (b) Satisfiability Don't-care
6. **(10 points)** What techniques are available for two-level logic minimization? Briefly (a word or two for each), what are the advantages and disadvantages of each?
7. **(5 points)** Cube $f(a, b) = a$. Cube $g(a, b) = a\bar{b}$. Does f cover g ? Does g cover f ?
8. **(5 points)** In two or three sentences, describe how Espresso searches for good implicants? Does it find all primes?
9. **(5 points)** Cofactor $f(a, b, c, d) = \bar{a}b + ac + ab\bar{d}$ by the cube $a\bar{b}$.
10. **(10 points)** Consider the following set of cubes:

$$\begin{array}{cccc|c}
 0 & 0 & X & 0 & 1 \\
 X & 1 & 0 & X & 1 \\
 1 & X & 0 & 1 & 1 \\
 1 & 1 & 1 & X & 1
 \end{array}$$

You would like to expand the 111X|1 cube to 11XX|1. However, you need to confirm that this expansion is valid.

- (a) Write down new cube that must be checked for containment.
- (b) Write down the set of cubes within which it must be contained.
- (c) Cofactor the set of cubes by the new cube and use the result to determine whether the cube is contained.
- (d) Repeatedly cofactor on unate variables of the cover until it is clear that the expansion is valid or invalid.
- (e) Explain why cofactoring on a unate variables (instead of non-unate variables) makes validity checking faster.